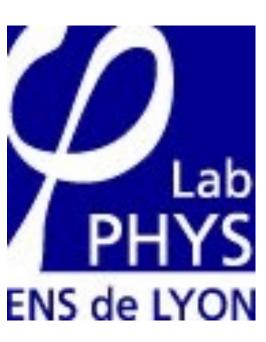


Graphs for data science and ML

Machine Learning for graphs and with graphs





Exploit the properties of the matrices of graphsSecond: find clusters, cut the graph

Cuts, clustering and communities

The good, the bad and the ugly

- Networks are often inhomogeneous, with important links, hubs, clusters, or communities (modules)
- These are observed in various types of data on networks: social, technological, biological,...
- Importance of cuts: the min-cut max-flow theorem.

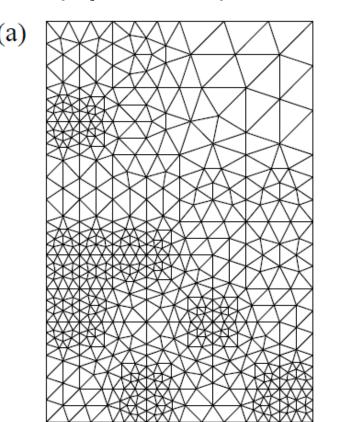
 These are two primal-dual linear programs.

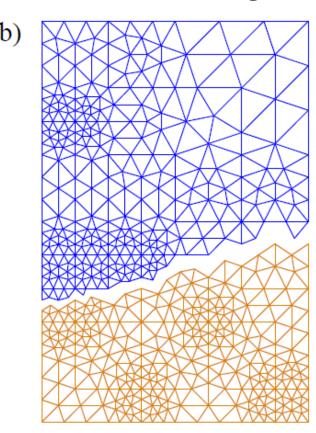
 The max value of a flow = the min capacity over all cuts.
- For clusters and communities, see the extensive surveys:

[S. Fortunato, *Physic Reports*, 2010]

[von Luxburg, Statistics and Computating, 2007]

• Example of (spectral) bisection on an irregular mesh





Exploit the properties of the matrices of graphs

Second: find clusters, cut the graph

Cuts, clustering and communities

The good, the bad and the ugly

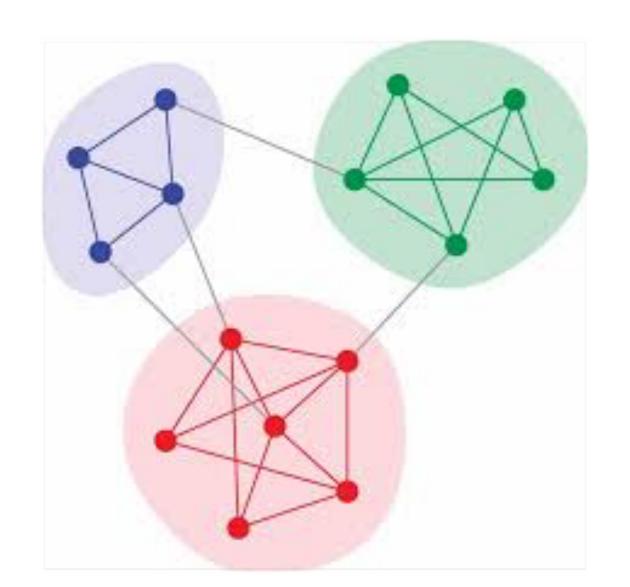
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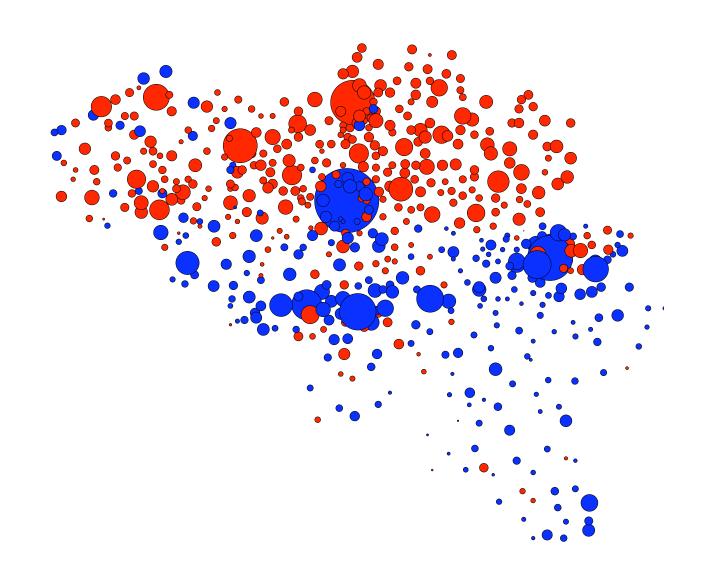
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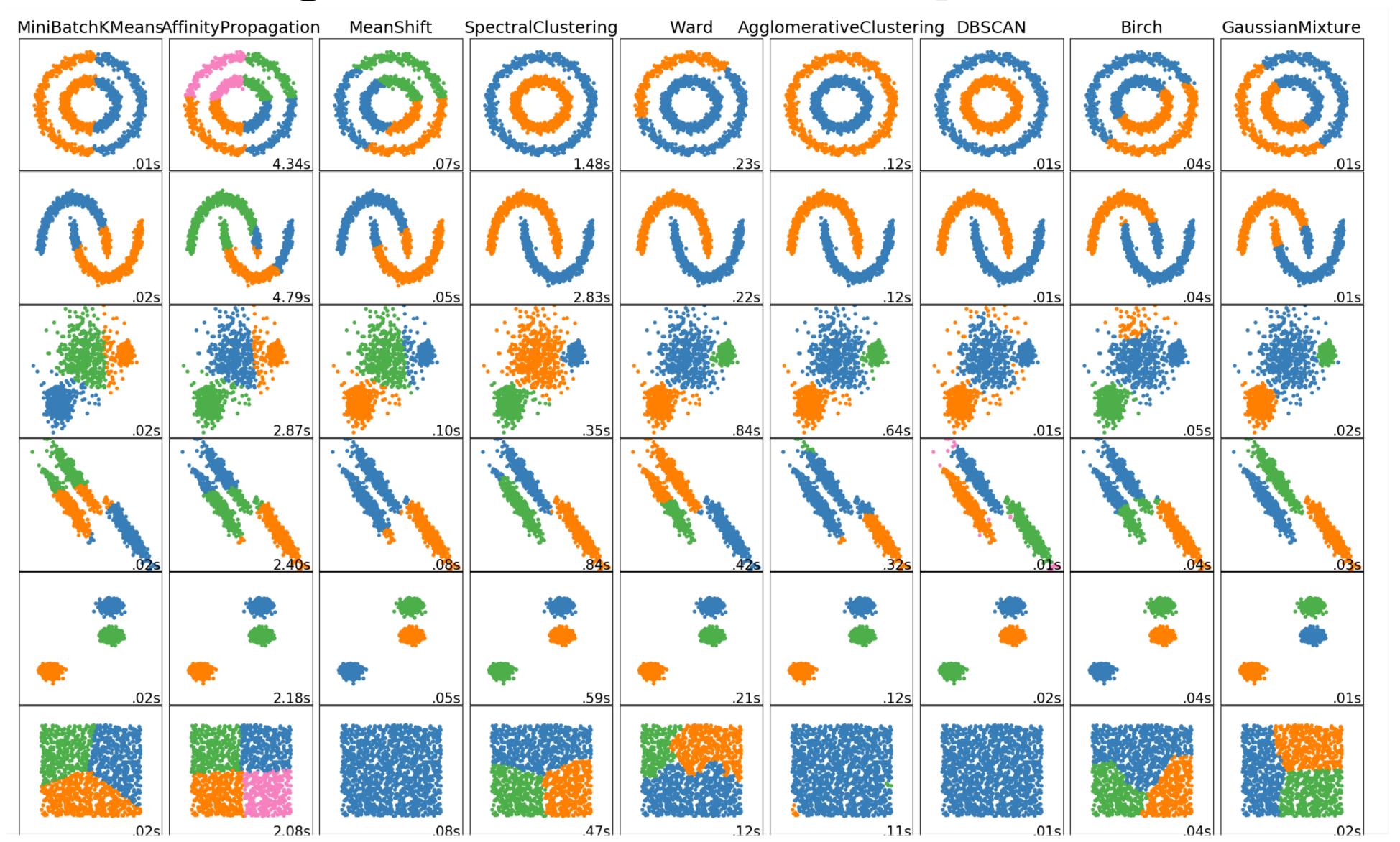
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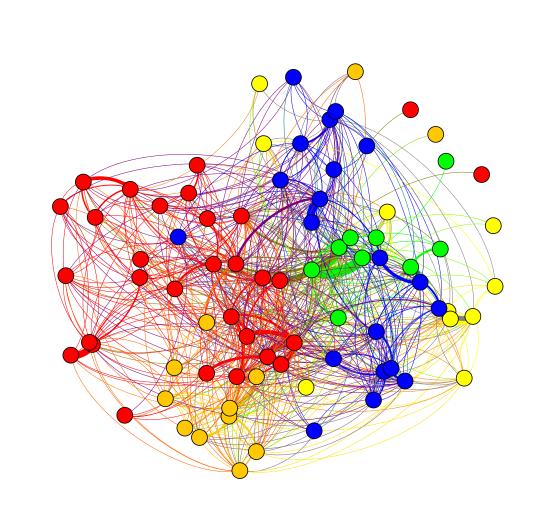


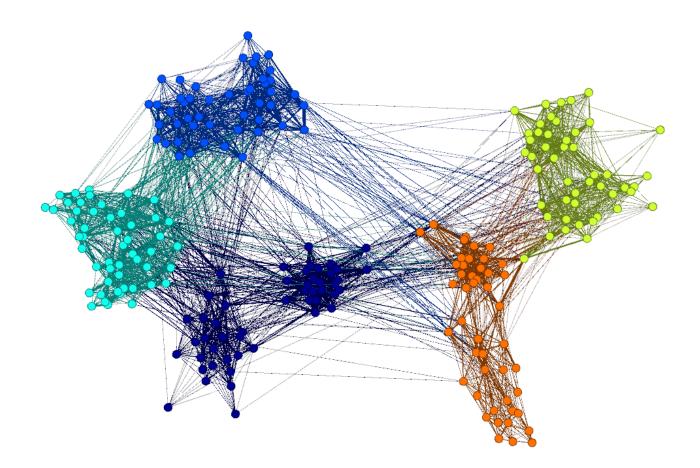


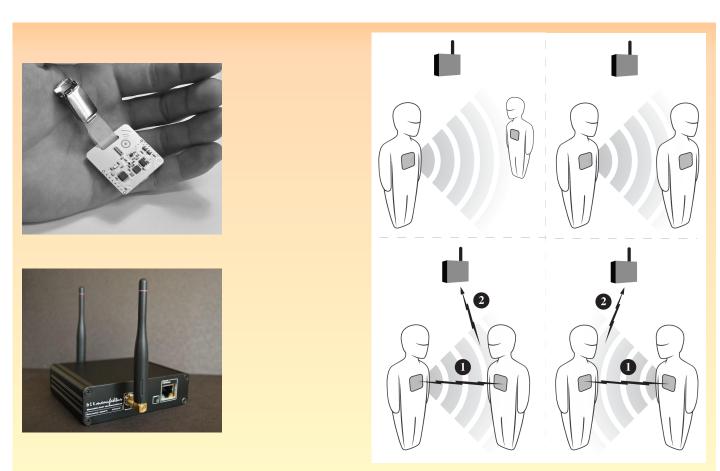
Clustering: a well known topic in data analysis



Social face-to-face interaction networks



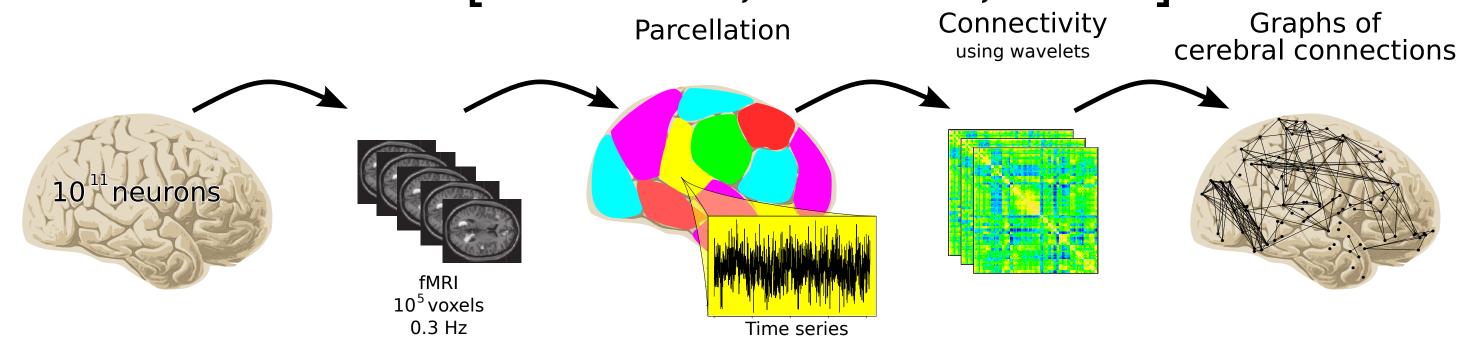




(Lab. physique, ENSL, 2013)

(école primaire, Sociopatterns)

Brain networks [Bullmore, Achard, 2006]



GRAPHSIP project challenges

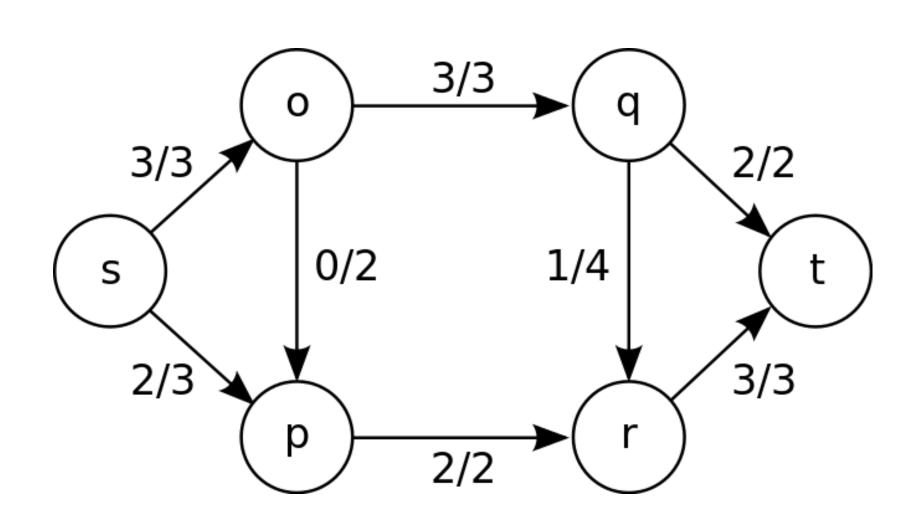
Clustering of graphs: also well known in Graph Theory

Min-Cut and Max-Flow

Graph cuts

- Graph cuts in 2 (or bisection): find the partition in two groups of nodes that minimize the cut size (i.e., the number of links cut)
- Exhaustive search can be computationally challenging
- About the problem of cuts:
 An important result is the min-cut max-flow theorem.

Min-cut pb and Max-flow pb are two primal-dual problems
The max value of a flow = the min capacity over all cuts
One possible solution from linear program



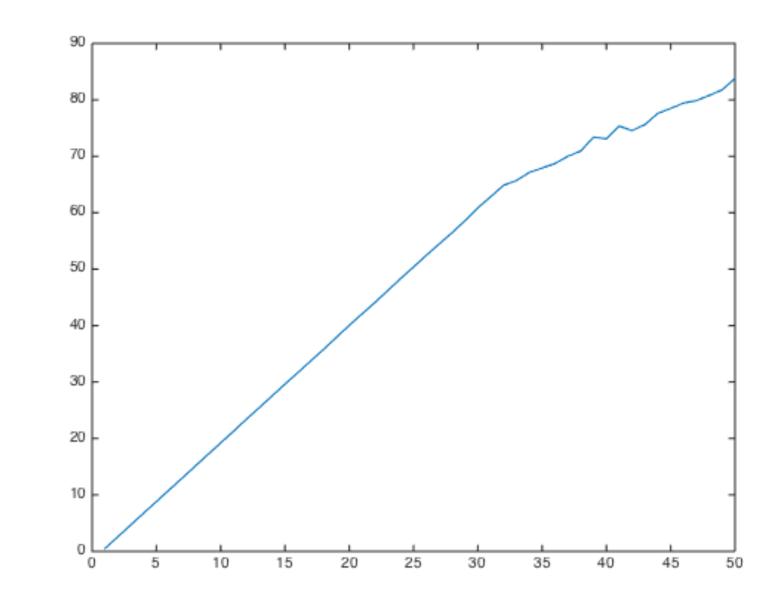
Some theoretical properties: Algebraic Connectivity

Multiplicity of eigenvalue 0 gives connectedness of graph

What if
$$\lambda_2 > 0$$
?

Experiment:

Gradually increase connections between two Erdos-Renyi subgraphs



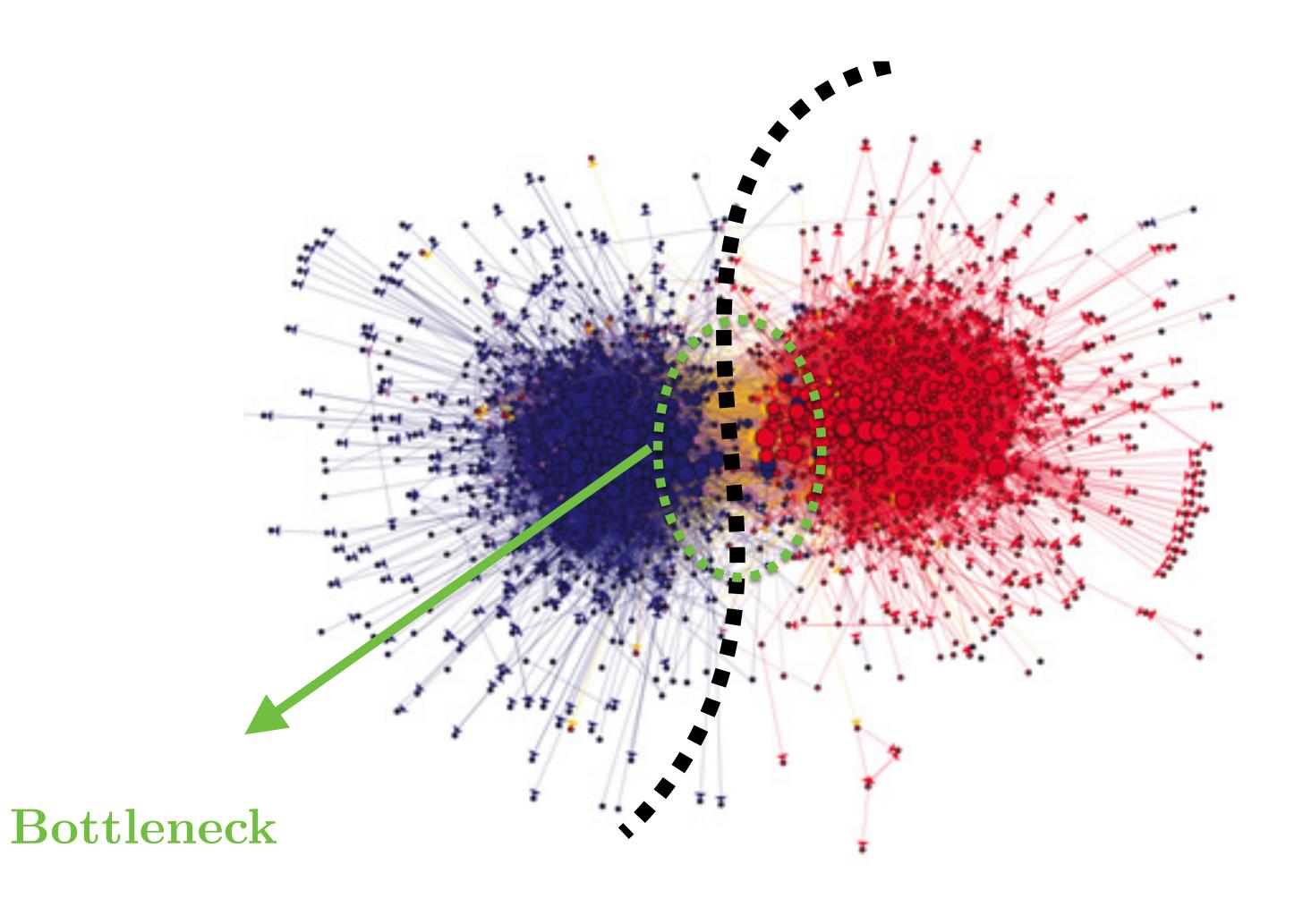
$$\lambda_2 \geqslant \frac{1}{\operatorname{vol}(G)d(G)}$$
 where $d(G)$ is the diameter of the graph

Some theoretical properties: The Cheeger Constant

Cheeger constant measures presence of a "bottleneck"

$$A \subset V \qquad \partial A = \left\{ (u, v) \in E \text{ s.t. } u \in A, v \in \overline{A} \right\}$$
$$\operatorname{vol}(A) = \sum_{u \in A} d(u)$$
$$h(G) = \min_{A \subset V} \left\{ \frac{|\partial A|}{\min\left(\operatorname{vol}(A), \operatorname{vol}(\overline{A})\right)} \text{ s.t. } 0 < |A| < \frac{1}{2}|V| \right\}$$

Some theoretical properties: The Cheeger Constant



Some theoretical properties: The Cheeger Constant

The Cheeger constant and algebraic connectivity are related by Cheeger inequalities. A simple example:

Theorem: Cheeger Inequality[Polya, Szego]

For a general graph G,

$$2h(G) \geqslant \lambda_2 \geqslant \frac{h^2(G)}{2}$$

Remark: the eigenvector associated to the algebraic connectivity is called the Fiedler vector

Some theoretical properties: The Fiedler Vector

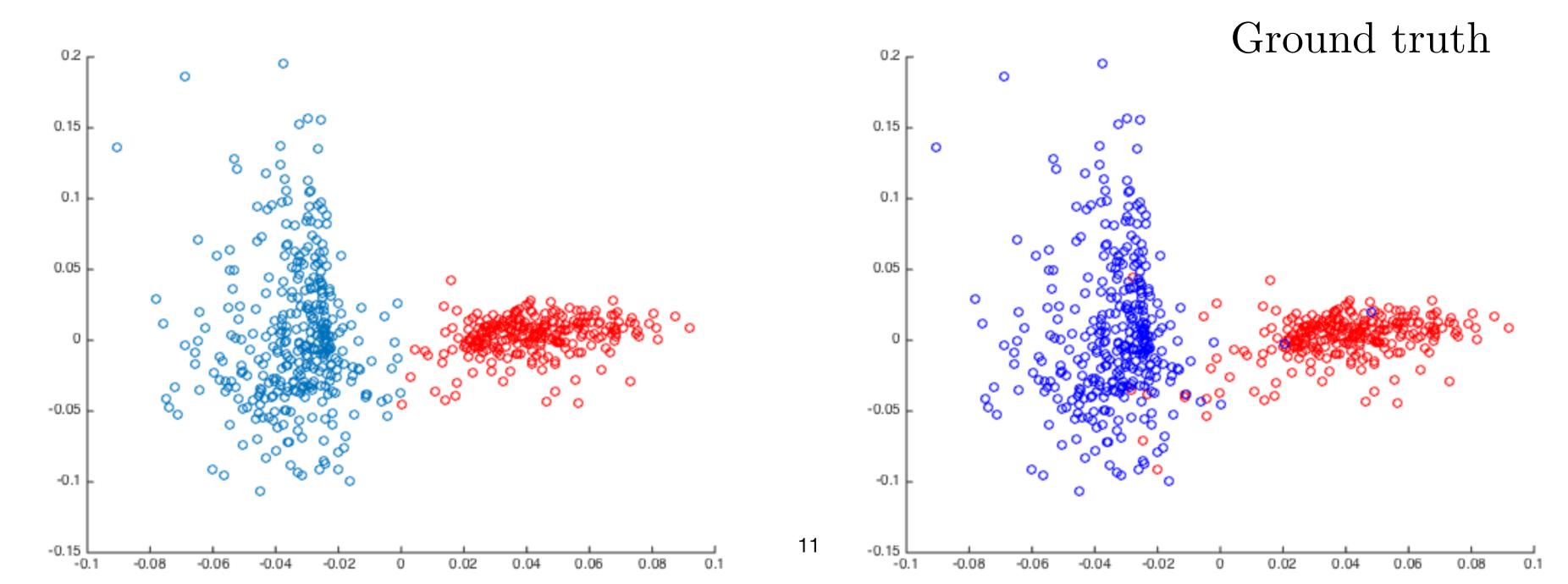
Set of 1490 US political blogs, labelled "Dem" or "Rep"

Hyperlinks among blogs

Removed small degrees (<12), keep N = 622 vertices

Compute normalised Laplacian, Fiedler vector

Assign attributes +1, -1 by sign of Fiedler vector



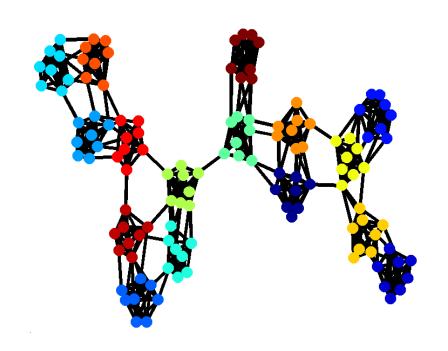
Clustering of graphs or Communities

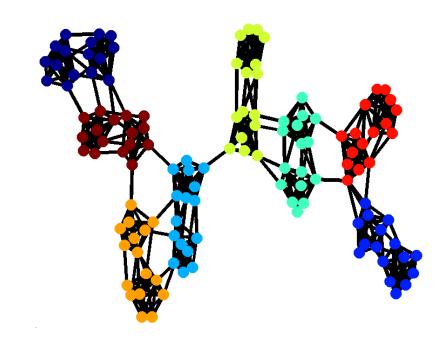
and their relation to the Laplacian eigenvectors

Example: graph with multiscale communities

finest scale (16 com.):

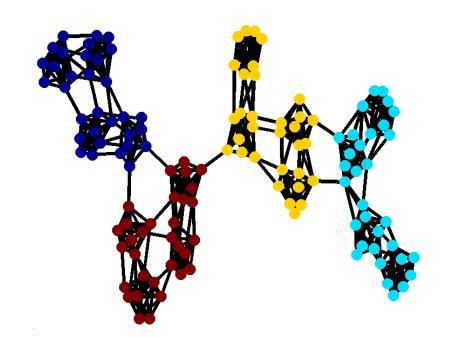


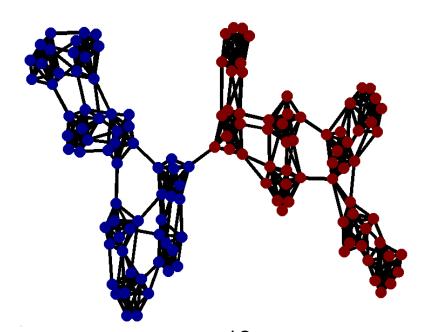




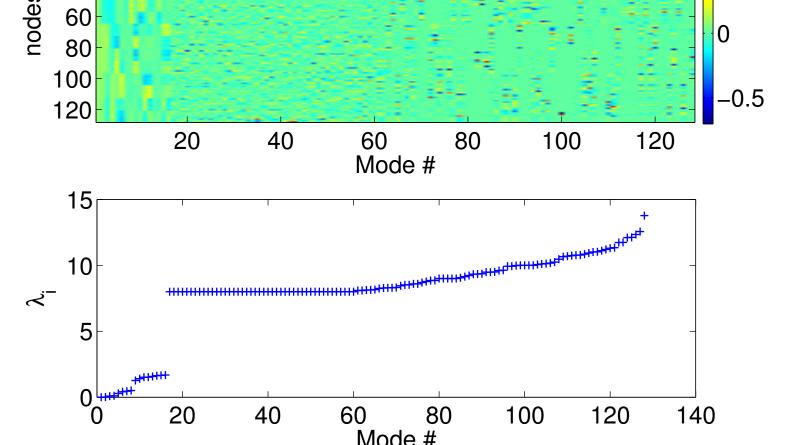
coarser scale (4 com.):

coarsest scale (2 com.):





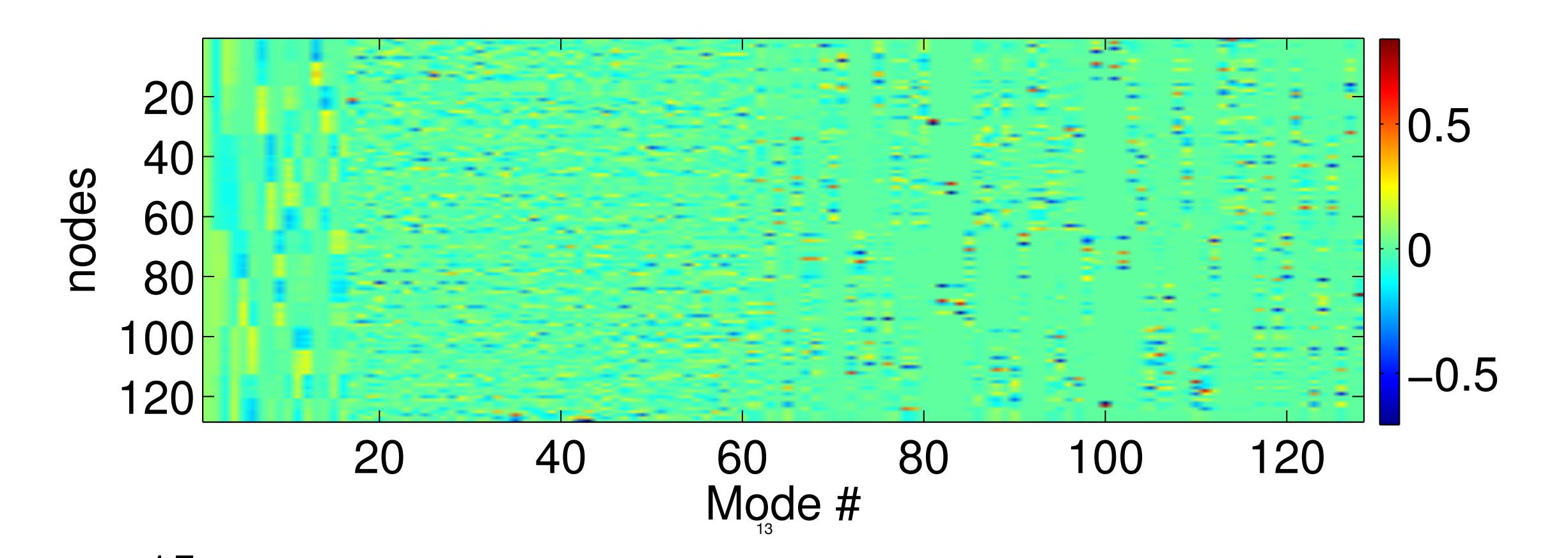
Spectral analysis: the χ_i and λ_i of this multiscale graph



Clustering of graphs or Communities

and their relation to the Laplacian eigenvectors

Spectral analysis: the χ_i and λ_i of this multiscale graph



U. von Luxburg, "A Tutorial on spectral clustering", Stat. Comput., 2007

- To cut a graph, one has to measure the size of the cut = # of cut edges
- then associate a cost function inspired by the Cheeger constant

$$C(A,B) := \sum_{i \in A, j \in B} \mathbf{W}[i,j]$$

$$RatioCut(A,\overline{A}) := \frac{1}{2} \frac{C(A,\overline{A})}{|A|} + \frac{1}{2} \frac{C(A,\overline{A})}{|\overline{A}|}$$

NormalizedCut
$$(A, \overline{A}) = \frac{1}{2} \frac{C(A, \overline{A})}{\text{vol}(A)} + \frac{1}{2} \frac{C(A, \overline{A})}{\text{vol}(\overline{A})}$$

How to minimise RatioCut? (a combinatorial problem)

• Re-write the problem with features indicative of each set:

$$f(i) = +\sqrt{\frac{|\bar{A}|}{|A|}} \text{ if } i \in A \quad \text{and } f(i) = -\sqrt{\frac{|\bar{A}|}{|A|}} \text{ if } i \in \bar{A}$$

- then $||f||_2 = \sqrt{|V|}$ and $f^\mathsf{T} \mathbf{1} = 0$
- One can compute the RatioCut : $f^{\mathsf{T}}\mathbf{L}f = |V|$ · RatioCut (A, \bar{A}) .
- The problem can be written as: $\min_{\bf f} \ {\bf f}^\top L \, {\bf f}$ such that ${\bf f}^\top {\bf 1} = 0, \ ||{\bf f}||_2 = \sqrt{|V|}$
 - With sign(f) an indicator function! <= still combinatorics

How to minimise RatioCut? (a relaxed, spectral, problem)

- The exact problem with f indicator function of A is (NP-)hard
- The same problem with any f is a relaxed version: looking for a smooth partition function

$$\operatorname{arg\,min}_{f} f^{T} \mathbf{L} f \text{ subject to } ||f|| = \sqrt{N}, \quad \langle f, 1 \rangle = 0$$

Solution (G connected): eigenvector of λ_2

Warning: recover partition after thresholding $f = sign(u_2)$

So we are back to the Fiedler vector!!!

RatioCut: generalization to k > 2

For more than two components, we look for a set of partition functions

$$F \in \mathbb{R}^{N \times k}$$
 $F[i,j] = f_j[i] = \begin{cases} 1/\sqrt{|A_j|} & \text{if } v_i \in A_j \\ 0 & \text{otherwise} \end{cases}$

Observe:
$$f_j^T \mathbf{L} f_j = \frac{\operatorname{Cut}(A_j, A_j)}{|A_j|}$$
 $F^T F = \mathbb{I}$

RatioCut
$$(A_1, \ldots, A_k) = \text{Tr}(F^T \mathbf{L} F)$$

Suggests the relaxed problem:

$$\arg\min_{F\in\mathbb{R}^{N\times k}}\operatorname{Tr}(F^T\mathbf{L}F)$$
 such that $F^TF=\mathbb{I}$

Unnormalized Spectral Clustering

This form of relaxed RatioCut = Unnormalized Spectral Clustering

$$\arg\min_{F\in\mathbb{R}^{N\times k}}\operatorname{Tr}(F^T\mathbf{L}F)$$
 such that $F^TF=\mathbb{I}$

Algorithm: Unnormalized Spectral Clustering

Compute the matrix F of first k eigenvectors of L Apply k-means to rows of F to obtain cluster assignments

Normalized Spectral Clustering

consider Normalized Cut, k=2

NormalizedCut
$$(A, \overline{A}) = \frac{1}{2} \frac{C(A, \overline{A})}{\text{vol}(A)} + \frac{1}{2} \frac{C(A, \overline{A})}{\text{vol}(\overline{A})}$$

Then:

$$f[i] = \begin{cases} \sqrt{\operatorname{vol}(\overline{A})/\operatorname{vol}(A)} & \text{if } v_i \in A\\ -\sqrt{\operatorname{vol}(A)/\operatorname{vol}(\overline{A})} & \text{otherwise} \end{cases}$$

Check that: $\langle \mathbf{D}f, 1 \rangle = 0$ $f^T \mathbf{D}f = \text{vol}(G)$

$$f^T \mathbf{L} f = \text{vol}(V) \text{NormalizedCut}(A, \overline{A})$$

$$\operatorname{arg\,min} f^T \mathbf{L} f \text{ subject to } f^T \mathbf{D} f = \operatorname{vol}(G), \quad \langle \mathbf{D} f, 1 \rangle = 0$$

$$g = \mathbf{D}^{1/2} f$$

 $\arg\min_{a} g^T \mathbf{L}_{\text{norm}}^{\mathbf{r}} g \text{ subject to } ||g||^2 = \text{vol}(G), \quad \langle g, \mathbf{D}^{1/2} 1 \rangle = 0$

Normalized Spectral Clustering

consider Normalized Cut, k>2

Then

$$F[i,j] = f_j[i] = \begin{cases} 1/\sqrt{\operatorname{vol}(A_j)} & \text{if } v_i \in A_j \\ 0 & \text{otherwise} \end{cases}$$

$$f_j^T \mathbf{L} f_j = rac{\operatorname{Cut}(A_j, \overline{A_j})}{\operatorname{vol}(A_j)}$$

$$F^TF = \mathbb{I}$$

$$F^T F = \mathbb{I} \qquad \qquad f_j^T \mathbf{D} f_j = 1$$

$$\arg\min_{H\in\mathbb{R}^{N\times k}}\operatorname{Tr}(H^T\mathbf{L}_{\mathrm{norm}}H)$$
 such that $H^TH=\mathbb{I}$

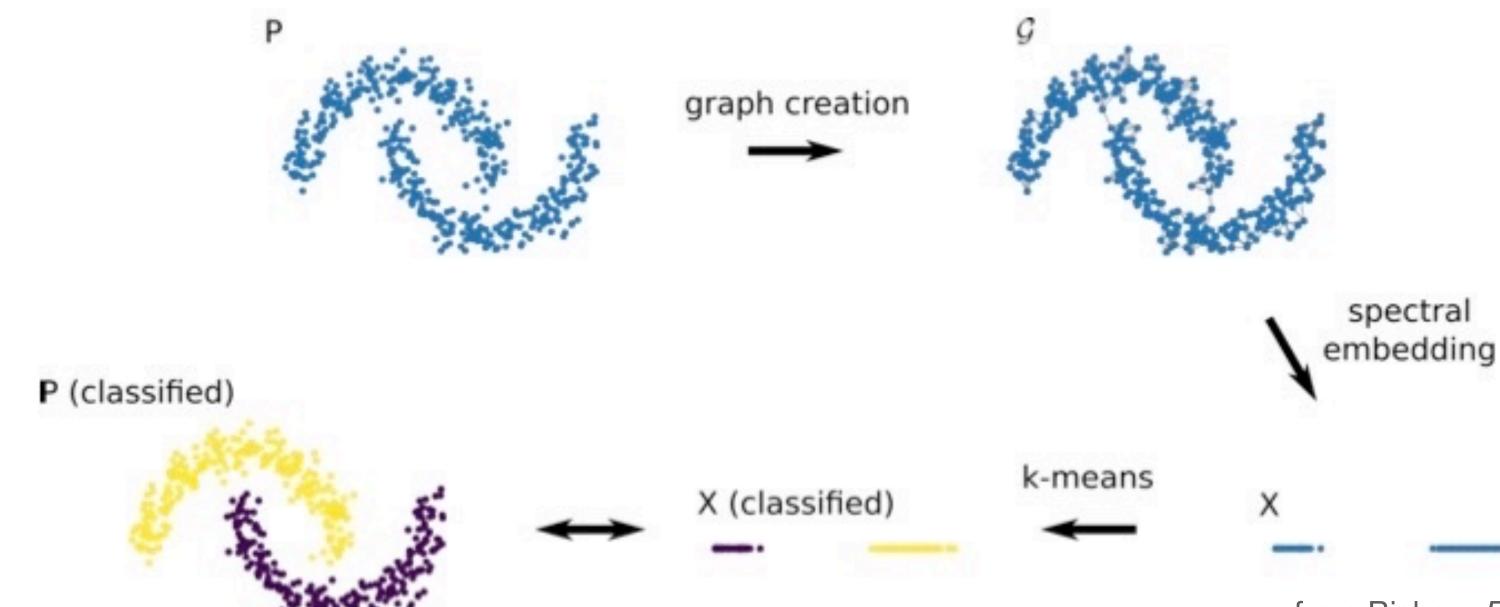
$$H = \mathbf{D}^{1/2} F$$

Algorithm: Normalized Spectral Clustering

Compute the matrix H of first k eigenvectors of \mathbf{L}_{norm} Apply k-means to rows of H to obtain cluster assignments

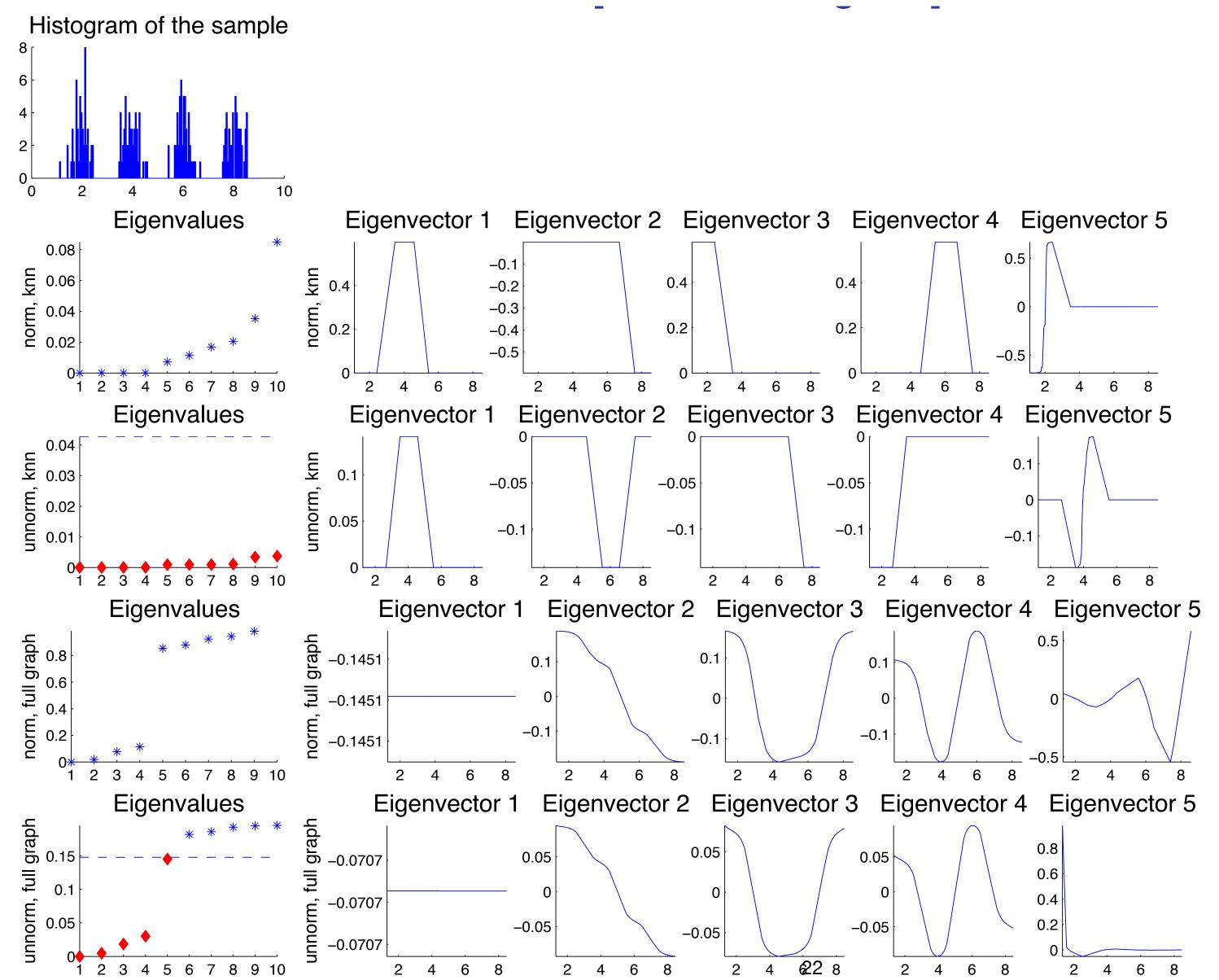
Spectral Clustering in a nutshell

- 1. Graph construction. A sparse similarity graph is built between the n points.
- Spectral embedding. The first k eigenvectors of a graph representative matrix (such as the Laplacian) are computed.
- 3. Clustering. k-means is performed on these spectral features, to obtain k clusters.



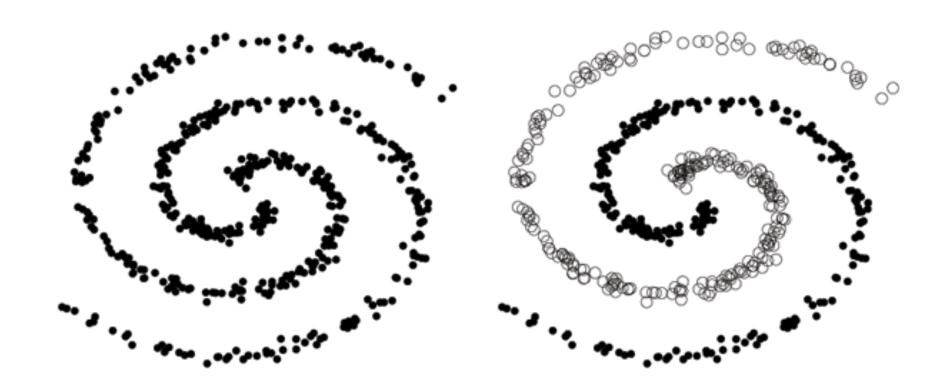
from Bishop, *Pattern Recognition and Machine Learning*, 2006 (chapt. 6)

Spectral Clustering in action



Spectral Clustering in action

In practice normalised spectral clustering is often preferred

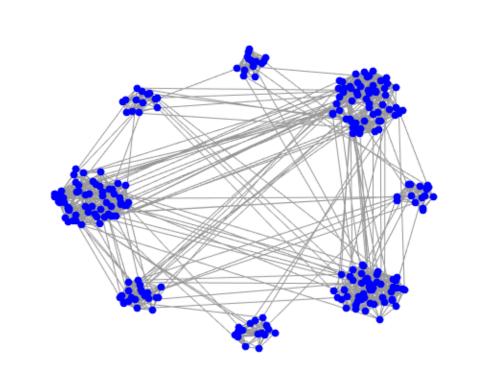


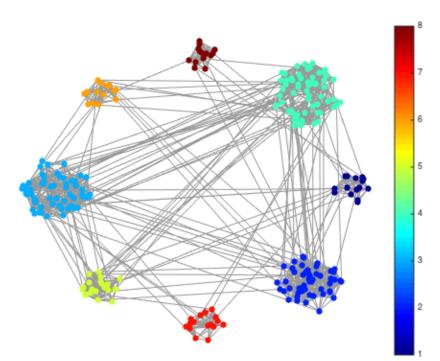
In practice the eigenvectors are "re-normalized" by the degrees $F = \mathbf{D}^{-1/2}H$ before k-means, because these are real cluster assignments

Rem: this is equivalent to using the "random walk Laplacian"

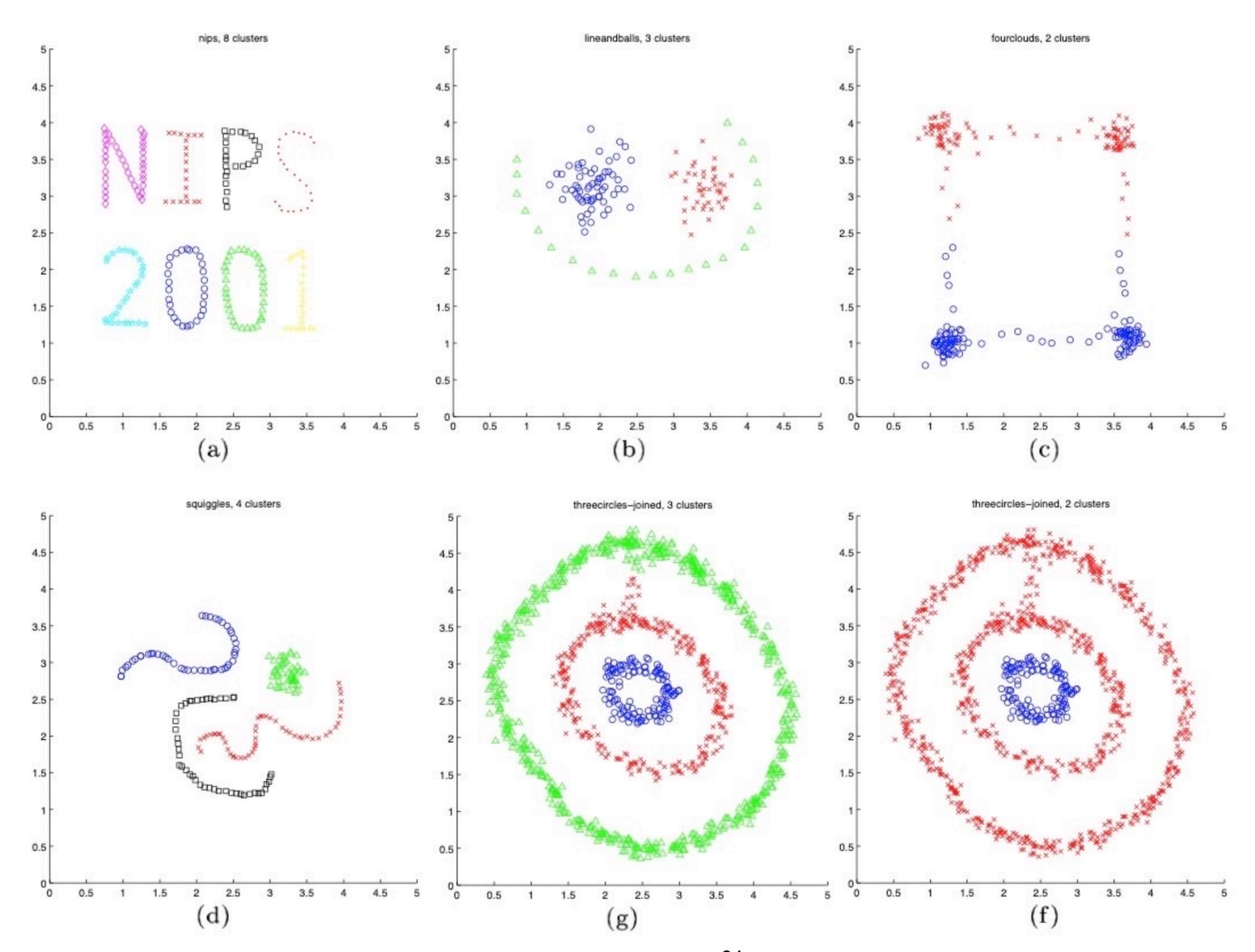
$$\mathbf{L}_{\mathrm{rw}} = \mathbf{D}^{-1} \mathbf{L}$$

If data has k clear clusters, there will be a gap in the Laplacian spectrum after the k-th eigenvalue. Use to choose k.





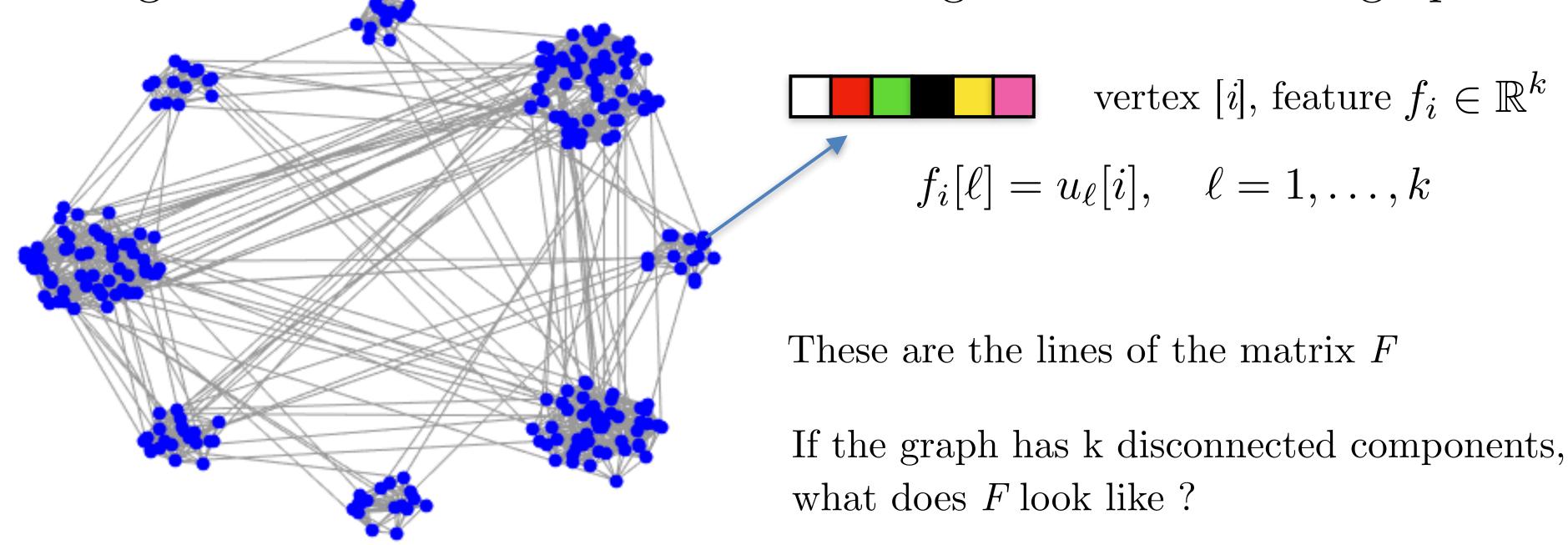
Spectral Clustering in action



Ng, Jordan, Weiss, NIPS'01

Spectral Clustering interpretation 1

At each vertex the algorithm associates a feature vector that represents the fine and large scale structure of that vertex's neighbourhood in the graph



k-means is then applied to these vectors to cluster into k clusters

In short, we transform the graph into a feature matrix and partition it.

Spectral Clustering interpretation 2

We are looking for k "partition signals (functions)"

$$f_{\ell}:V\mapsto\mathbb{R}$$

In the ideal case (k disconnected components)

$$f_{\ell}[i] = \begin{cases} c_i & \text{if } i \in \text{cluster } \ell \\ 0 & \text{otherwise} \end{cases}$$

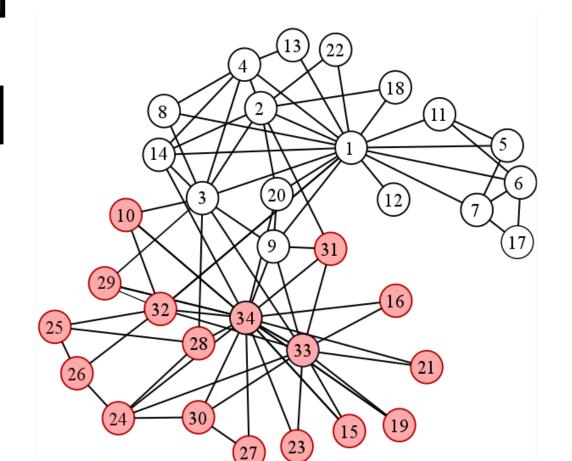
7 6 5 4 4 3 2 2

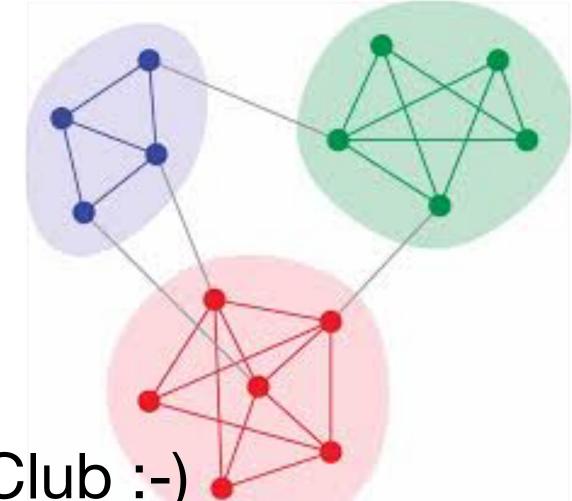
These are maximally smooth graph signals: $f_{\ell}^T \mathbf{L} f_{\ell} = 0$

Exploit the properties of the matrices of graphs

Third: find communities in complex networks

- Communities are more loosely defined:
 - usually = nodes more connected together than with the outside
- Many methods (thousands of papers):
 - Modularity [Girvan, Newman, 2004]
 - Infomap [Rosvall, Bergstrom, 2009]
 - Stochastic Block Models





Win a prize at NetSci: be 1st to talk about the Zachary Karate Club:-)

Summary up to now:

Exploit the spectral analysis of graphs

- Graph spectral analysis: eigenvectors and eigenvalues of matrices $A,\,L,\,L_{rw},\dots$
- They have interesting properties, especially for L: new basis of representation, analog for Fourier oscillations, study of smoothness on graphs,...
- Centralities: can be studied with these matrices
- Clusters: can be found with them as well (in a relaxed way)

More processing methods to come!